

Join us at the ...

**11th Conference of the Fachgruppe Computeralgebra**

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- **Young Talent Award** for best contributed talk: € 500 plus invitation as main speaker for next conference



**Fachgruppe**  
Computeralgebra



*Die Fachgruppe Computeralgebra fördert Lehre, Forschung, Entwicklung, Anwendungen, Informationsaustausch und Zusammenarbeit auf dem Gebiet der Computeralgebra in Deutschland.*

[↗ fachgruppe-computeralgebra.de](https://fachgruppe-computeralgebra.de)

Fachgruppe der GI in Kooperation mit der DMV und GAMM



Gesellschaft für Informatik



Deutsche Mathematiker-Vereinigung



Gesellschaft für Angewandte Mathematik und Mechanik

# The state of OSCAR

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John Abbott, Thomas Breuer, Lars Göttgens, Max Horn



March 26, 2025

The state of OSCAR

*Max Horn*

S-characters

*Thomas Breuer*

Algebraic Lie Theory

*Lars Göttgens*

Practical Fast Linear Algebra Over  $\mathbb{Z}$

*John Abbott*

Summing it up

*Max Horn*

# The state of OSCAR

Max Horn



## What happened so far (abridged version)

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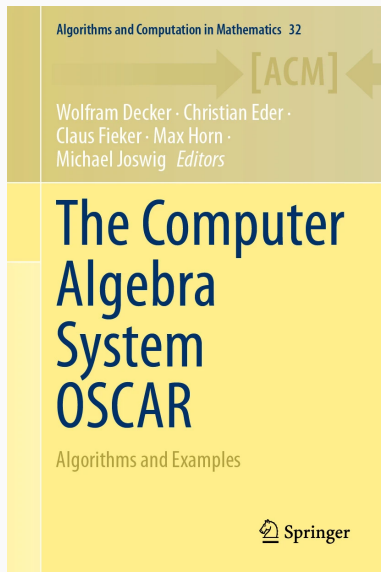
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- 02/2025: relaunch of  [oscar-system.org](https://oscar-system.org)
- 03/2025: OSCAR v1.3.1 released
- In the past 12 months, 47 people contributed code to our main code repository

# The OSCAR book

The OSCAR book is out!

Get it while it's hot

See  [book.oscar-system.org](http://book.oscar-system.org)



# And now to something different

The state of OSCAR

*Max Horn*

S-characters

*Thomas Breuer*

Algebraic Lie Theory

*Lars Göttgens*

Practical Fast Linear Algebra Over  $\mathbb{Z}$

*John Abbott*

Summing it up

*Max Horn*

# S-characters

Thomas Breuer

---

- joint work with Michael Joswig and Gunter Malle
- <https://arxiv.org/abs/2408.16785>
- <https://github.com/oscar-system/FriendsOfOscar/tree/main/Zeros%2Bof%2BS-Characters>

## S-characters: Definitions

$G$  finite group

$g_1 = 1, g_2, \dots, g_n$  conjugacy class representatives

$\text{Irr}(G) = \{\chi_1 = 1_G, \chi_2, \dots, \chi_n\}$  irreducible characters

$A = [\chi_i(g_j)]_{i,j} \in \mathbb{C}^{n \times n}$  character table

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$\text{SCoord}(G) = \{x \in \mathbb{Z}^n; x \cdot A \in (\mathbb{R}^n)^{\geq 0}, x_1 = 1\}$

$\text{SChar}(G) = \{\sum_{i=1}^n x_i \chi_i; x \in \text{SCoord}(G)\}$  S-characters

(Zhمود' 1995)

# Example: S-characters of $A_5$

$$G = A_5$$

1A 2A 3A 5A 5B

1	1	1	1	1
3	-1	.	-b5	*
3	-1	.	*	-b5
4	.	1	-1	-1
5	1	-1	.	.

$$b5 = \frac{1}{2} (-1 + \sqrt{5})$$

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	1A	2A	3A	5A	5B
1	-1	-1	0	1	0	4	0	0	0
1	0	0	0	0	1	1	1	1	1
1	0	0	0	1	6	2	0	1	1
1	0	0	1	-1	0	0	3	0	0
1	0	0	1	0	5	1	2	0	0
1	0	0	1	1	10	2	1	0	0
1	0	0	1	2	15	3	0	0	0
1	0	1	-1	0	0	0	0	b5+3	-b5+2
1	0	1	0	0	4	0	1	b5+2	-b5+1
1	0	1	0	1	9	1	0	b5+2	-b5+1
1	1	0	-1	0	0	0	0	-b5+2	b5+3
1	1	0	0	0	4	0	1	-b5+1	b5+2
1	1	0	0	1	9	1	0	-b5+1	b5+2
1	1	1	0	1	12	0	0	2	2
1	1	1	1	1	16	0	1	1	1
1	1	1	1	2	21	1	0	1	1
1	1	1	2	1	20	0	2	0	0
1	1	1	2	2	25	1	1	0	0
1	1	1	2	3	30	2	0	0	0
1	1	2	1	2	24	0	0	b5+2	-b5+1
1	2	1	1	2	24	0	0	-b5+1	b5+2
1	2	2	2	3	36	0	0	1	1
1	2	2	3	3	40	0	1	0	0
1	2	2	3	4	45	1	0	0	0
1	3	3	4	5	60	0	0	0	0

## Zeros of $S$ -characters

- Each  $S$ -character except  $1_G$  has a zero.
- If  $H \leq G$  then  $1_H^G$  is an  $S$ -character.  
(Permutation character)
- If  $H < G$  then  $1_H^G$  is zero at an element of prime power order.  
(Fein-Kantor-Schacher 1981)

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Question:  $\text{SChar}^{\text{PPP}}(G) \neq \{1_G\}$ ?

## Compute $\text{SChar}^{\text{ppp}}(G)$ : Situation

$A = [\chi_i(g_j)]_{i,j}$  character table

Assume  $A \in \mathbb{R}^{n \times n}$ .

$$\text{SCoord}(G) = \{x \in \mathbb{Z}^n; x \cdot A \in (\mathbb{R}^n)^{\geq 0}, x_1 = 1\}$$

**Combinatorially:**

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( $G = A_5$ : 4851 candidates,  
 $G = A_8$ :  $\approx 10^{20}$  candidates)

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**Geometrically:**

$$S(G) = \{x \in \mathbb{R}^n; x \cdot A \in (\mathbb{R}^n)^{\geq 0}, x_1 = 1\}$$

is a bounded polytope,

$S\text{Coord}(G)$  are its lattice points.

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**Note:** All lattice points except  $1_G$  lie on the boundary.

## Compute $S\text{Char}^{PPP}(G)$ : Idea

- **Use OSCAR:**

Create the polytope  $S(G)$ ,  
enumerate its lattice points,  
check which of them belong to  $S\text{Char}^{PPP}(G)$ .

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The polytope is defined over an embedded number field.

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check which of them belong to  $S\text{Char}^{\text{PPP}}(G)$ .

- **Challenge:**

The polytope is defined over an embedded number field.

- **Improvements:**

Prescribe nonzero values a priori on rational columns.  
Use projections to factor groups where possible.

## Compute $S\text{Char}^{ppp}(A_8)$

```
julia> using S_characters, Oscar
julia> tbls = map(character_table, ["A8", "2.A8"]);
julia> @time length(s_characters(tbls[1]))
2.462459 seconds (27.24 M allocations: 416.007 MiB)
3635
julia> @time length(s_characters(tbls[1], ppow_nonzero = true))
0.081231 seconds (543.64 k allocations: 4.496 MiB)
1
```

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0.081231 seconds (543.64 k allocations: 4.496 MiB)  
1  
julia> @time length(s_characters(tbls[2], ppow_nonzero = true))  
0.308319 seconds (3.79 M allocations: 30.041 MiB)  
3  
julia> @time length(s_characters_via_factor_group(tbls[2],  
[1, 2], ppow_nonzero = true))  
0.186391 seconds (656.97 k allocations: 6.412 MiB)  
3
```

# Results: $\text{SChar}^{\text{PPP}}(G) \neq \{1_G\}$

$G$	#classes	#real	#rat.	#SChar <sup>PPP</sup>	#virt. SChar <sup>PPP</sup>
$A_8$	14	12	12	1	0
$2.A_8$	23	18	18	2	1
$2^6:A_8$	41	39	39	8	1
$A_9$	18	17	17	3	0
$2.A_9$	30	26	26	10	7
$A_{10}$	24	24	23	1	0
$2.A_{10}$	39	38	35	3	2
$A_{11}$	31	30	29	6	0
$2.A_{11}$	49	47	43	53	26
$M_{12}$	15	14	14	1	0
$M_{12}.2$	21	21	19	1	0
$M_{24}$	26	21	21	223	4
$J_1$	15	15	10	1	0
$J_2$	21	21	16	155	8
$2.J_2$	38	37	28	2571	1280
$J_2.2$	27	27	26	12	1
$J_3$	21	20	14	5	0
$M^cL$	24	18	18	2588	65
$HS$	24	22	22	93	2
$2.HS$	42	34	33	2211	1094

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$G$	#classes	#real	#rat.	# $\text{SChar}^{PPP}$	#virt. $\text{SChar}^{PPP}$
$L_4(3)$	29	26	25	6	0
$L_5(2)$	27	20	18	11	0
$2^5.L_5(2)$	41	29	27	2696	118
$S_4(4)$	27	27	18	133	10
$S_4(4).2$	30	29	28	5	1
$S_4(5)$	34	34	25	76	11
$S_6(2)$	30	30	30	1	1
$U_4(2)$	20	15	15	2	1
$2.U_4(2)$	34	24	24	1	1
${}^2F_4(2)'$	22	19	16	71	2
${}^2F_4(2)'.2$	29	23	23	7	3
${}^2G_2(27)$	35	32	13	9	0
$G_2(3)$	23	22	21	24	0

## Results: $\text{SChar}^{\text{PPP}}(G) \neq \{1_G\}$

Our examples of  $\psi \in \text{SChar}^{\text{PPP}}(G) \setminus \{1_G\}$ :

- All known  $\psi$  are rational-valued.  
(Are there non-rational examples?)
- $G = A_8$  is the smallest such Atlas group.
- Checked simple groups by increasing conjugacy class number:

The computation did not finish for

$O'_N, G_2(4), He, {}^3D_4(2), Ru, O_8^-(2), Co_3, A_{12}$

(with 128 GB RAM).

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- Checked for *rational*  $\psi$  in  $G$  with at most 40 rational classes:

No example  $\psi$  known for  $|G| \leq 1535$ .

No example  $\psi$  known for solvable  $G$ .

(Are there solvable examples?)

**We can show:**

Solvable groups do not have *proper characters* as examples.)

S-characters are ...

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... virtual characters and real.

S-characters are ...

... virtual characters and real.

Virtual Reality?

# Algebraic Lie Theory

## Lars Göttingen

---

## Feature overview

- (generalized) Cartan matrices,
- Root systems,
- Weyl groups,
- Lie algebras,
- Lie subalgebras, ideals
- Lie algebra modules
- Morphisms of Lie algebras, morphisms of modules
- Representation theory helpers ( $L$  semisimple,  $\text{char } K = 0$ ):
  - dimension formula, character formula, character formula for Demazure modules

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Non-trivial contributions by: A. Della Vecchia, LG, M. Horn, J. Nohilly, J. Peters, F. Röhrich, L. Voggesberger, T. Wiedemann

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- Dynkin type, or list thereof

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julia> rs = root_system([(G, 2), (E, 7)])  
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julia> rs = root_system([(:G, 2), (:E, 7)])  
Root system of rank 9  
of type G2 x E7
```

Running example:

```
julia> rs = root_system(:B, 3)  
Root system of rank 3  
of type B3
```

## Root systems: reflections at simple and arbitrary roots

```
julia> positive_roots(rs)
9-element Vector{RootSpaceElem}:
```

$\alpha_1$

$\alpha_2$

$\alpha_3$

$\alpha_1 + \alpha_2$

$\alpha_2 + \alpha_3$

$\alpha_2 + 2\alpha_3$

$\alpha_1 + \alpha_2 + \alpha_3$

$\alpha_1 + \alpha_2 + 2\alpha_3$

$\alpha_1 + 2\alpha_2 + 2\alpha_3$

```
julia> reflect(root(rs, 1), 2) #  $\sigma_2(\alpha_1)$ 
```

$\alpha_1 + \alpha_2$

```
julia> reflect(root(rs, 1), root(rs, 4)) #  $\sigma_{\{\alpha_1+\alpha_2\}}(\alpha_1)$ 
```

$-\alpha_2$

## Root systems: root arithmetics

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julia> positive_roots(rs)
9-element Vector{RootSpaceElem}:
```

$\alpha_1$

$\alpha_2$

$\alpha_3$

$\alpha_1 + \alpha_2$

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$\alpha_2 + 2\alpha_3$

$\alpha_1 + \alpha_2 + \alpha_3$

$\alpha_1 + \alpha_2 + 2\alpha_3$

$\alpha_1 + 2\alpha_2 + 2\alpha_3$

```
julia> is_positive_root_with_index(root(rs, 5) + root(rs, 3))
(true, 6)
```

```
julia> is_positive_root_with_index(root(rs, 5) + 2*root(rs, 2))
(false, 0)
```

# Weyl groups

Root systems and Weyl groups are intrinsically connected.

```
julia> W = weyl_group(rs)
Weyl group
  of root system of rank 3
  of type B3
```

```
julia> s1,s2,s3 = gens(W)
[s1, s2, s3]
```

```
julia> one(W)
id
```

```
julia> longest_element(W)
s1 * s2 * s1 * s3 * s2 * s1 * s3 * s2 * s3
```

## Weyl groups: element representations

```
julia> x = W([3,2,1])  
s3 * s2 * s1
```

```
julia> word(x)  
UInt8[0x03, 0x02, 0x01]
```

```
julia> letters(x)  
Int[3, 2, 1]
```

```
julia> syllables(x)  
3-element Vector{Pair{Int64, ZZRingElem}}:  
 3 => 1  
 2 => 1  
 1 => 1
```

All of these representations can also be passed to `W` to construct an element.

## Weyl groups: normal form

Elements are always stored as a word in *short-lex normal form*, i.e., the lexicographic smallest word of minimal length.

```
julia> W([1,2,2,3])  
s1 * s3
```

## Weyl groups: normal form

Elements are always stored as a word in *short-lex normal form*, i.e., the lexicographic smallest word of minimal length.

```
julia> W([1,2,2,3])  
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julia> W([3,2,1,2,3])  
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julia> W([1,2,2,3])  
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```
julia> W([3,2,1,2,3])  
s1 * s3 * s2 * s1 * s3
```

```
julia> x = s3 * s1 * s2 * s3 * s2 * s1  
s1 * s2 * s3 * s2 * s1 * s3
```

## Weyl groups: normal form

Elements are always stored as a word in *short-lex normal form*, i.e., the lexicographic smallest word of minimal length.

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```

```
julia> W([3,2,1,2,3])  
s1 * s3 * s2 * s1 * s3
```

```
julia> x = s3 * s1 * s2 * s3 * s2 * s1  
s1 * s2 * s3 * s2 * s1 * s3
```

```
julia> inv(x)  
s1 * s2 * s3 * s2 * s1 * s3
```

## Weyl groups: all reduced expressions of an element

```
julia> x = s3 * s1 * s2 * s3 * s2 * s1  
s1 * s2 * s3 * s2 * s1 * s3
```

```
julia> reduced_expressions(x)  
ReducedExpressionIterator(s1 * s2 * s3 * s2 * s1 * s3, false)
```

## Weyl groups: all reduced expressions of an element

```
julia> x = s3 * s1 * s2 * s3 * s2 * s1  
s1 * s2 * s3 * s2 * s1 * s3
```

```
julia> reduced_expressions(x)  
ReducedExpressionIterator(s1 * s2 * s3 * s2 * s1 * s3, false)
```

```
julia> reduced_expressions(x) |> collect .|> Vector{Int}
```

```
4-element Vector{Vector{Int64}}:
```

```
[1, 2, 3, 2, 1, 3]
```

```
[1, 2, 3, 2, 3, 1]
```

```
[1, 3, 2, 3, 2, 1]
```

```
[3, 1, 2, 3, 2, 1]
```

## Weyl groups: all reduced expressions of an element

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julia> x = s3 * s1 * s2 * s3 * s2 * s1  
s1 * s2 * s3 * s2 * s1 * s3
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ReducedExpressionIterator(s1 * s2 * s3 * s2 * s1 * s3, false)
```

```
julia> reduced_expressions(x) |> collect .|> Vector{Int}  
4-element Vector{Vector{Int64}}:  
 [1, 2, 3, 2, 1, 3]  
 [1, 2, 3, 2, 3, 1]  
 [1, 3, 2, 3, 2, 1]  
 [3, 1, 2, 3, 2, 1]
```

```
julia> reduced_expressions(x; up_to_commutation=true) |>  
↪ collect .|> Vector{Int}  
4-element Vector{Vector{Int64}}:  
 [1, 2, 3, 2, 1, 3]  
 [1, 3, 2, 3, 2, 1]
```

# Weyl groups: operation on roots and weights

Warning: **RIGHT** action!

## Weyl groups: operation on roots and weights

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```
julia> a = positive_root(rs, 4)  
 $\alpha_1 + \alpha_2$ 
```

```
julia> a * s1  
 $\alpha_2$ 
```

## Weyl groups: operation on roots and weights

Warning: **RIGHT** action!

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julia> a = positive_root(rs, 4)
 $\alpha_1 + \alpha_2$ 
```

```
julia> a * s1
 $\alpha_2$ 
```

```
julia> w = fundamental_weight(rs, 1)
 $\omega_1$ 
```

```
julia> w * s1
 $-\omega_1 + \omega_2$ 
```

## Weyl groups: conversion to other group types

Challenge: Weyl groups are the first (widely-used, multiplicative) groups in OSCAR, that are not backed by GAP.

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(temporary) solution: convert to a group of another type

```
julia> iso = isomorphism(PermGroup, W)
```

```
Map defined by a julia-function with inverse  
  from Weyl group of root system of type B3  
  to permutation group of degree 6 and order 48
```

```
julia> FPGroup(W)
```

```
Finitely presented group of order 48
```

```
julia> geometric_representation(W)
```

```
(Matrix group of degree 3 over ZZ, Map: W -> matrix group)
```

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  to permutation group of degree 6 and order 48
```

```
julia> FPGGroup(W)
```

```
Finitely presented group of order 48
```

```
julia> geometric_representation(W)
```

```
(Matrix group of degree 3 over ZZ, Map: W -> matrix group)
```

A lot of ongoing effort by many people to make Weyl groups more usable, and add more features.

## Lie algebras: construction

Fix a field  $F$ .

Construction via:

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- structure constant table (as `Array{elem_type(F), 3}` or `Matrix{sparse_row_type(F)}`)

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- structure constant table (as `Array{elem_type(F), 3}` or `Matrix{sparse_row_type(F)}`)
- root system or Dynkin type

Running example:

```
julia> sc = fill(sparse_row(QQ), 14, 14);
julia> sc[1,3] = sparse_row(QQ, [3,8,10,11,13],
↪ [3,2,-4,-2,2]);
julia> sc[1,4] = sparse_row(QQ, [4,8,9,10,11,13],
↪ [5,4//3,-3,-8//3,8//3,-8//3]);
# ...
julia> L = lie_algebra(QQ, sc, ["b_$(i)" for i in 1:14])
Abstract Lie algebra
  of dimension 14
```

# Lie algebras: properties

```
julia> dim(L)
```

```
14
```

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```
julia> dim(L)
```

```
14
```

```
julia> is_nilpotent(L)
```

```
false
```

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```
julia> dim(L)
```

```
14
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```
julia> is_nilpotent(L)
```

```
false
```

```
julia> is_solvable(L)
```

```
false
```

# Lie algebras: properties

```
julia> dim(L)
```

```
14
```

```
julia> is_nilpotent(L)
```

```
false
```

```
julia> is_solvable(L)
```

```
false
```

```
julia> is_perfect(L)
```

```
true
```

# Lie algebras: properties

```
julia> dim(L)
```

```
14
```

```
julia> is_nilpotent(L)
```

```
false
```

```
julia> is_solvable(L)
```

```
false
```

```
julia> is_perfect(L)
```

```
true
```

```
julia> rank(killing_matrix(L))
```

```
14
```

# Lie algebras: properties

```
julia> dim(L)
```

```
14
```

```
julia> is_nilpotent(L)
```

```
false
```

```
julia> is_solvable(L)
```

```
false
```

```
julia> is_perfect(L)
```

```
true
```

```
julia> rank(killing_matrix(L))
```

```
14
```

```
julia> is_semisimple(L)
```

```
true
```

## Lie algebras: Cartan subalgebra

```
julia> csa = cartan_subalgebra(L)
Lie subalgebra
  of dimension 2
of abstract Lie algebra over QQ
```

## Lie algebras: Cartan subalgebra

```
julia> csa = cartan_subalgebra(L)
Lie subalgebra
  of dimension 2
of abstract Lie algebra over QQ

julia> basis(csa)
2-element Vector{AbstractLieAlgebraElem{QQFieldElem}}:
 b_1
 b_2
```

## Lie algebras: Cartan subalgebra

```
julia> csa = cartan_subalgebra(L)
Lie subalgebra
  of dimension 2
of abstract Lie algebra over QQ

julia> basis(csa)
2-element Vector{AbstractLieAlgebraElem{QQFieldElem}}:
 b_1
 b_2

julia> normalizer(L, csa) == csa
true

julia> is_nilpotent(lie_algebra(csa)[1])
true
```

## Lie algebras: root system

```
julia> rs = root_system(L)
Root system of rank 2
of type G2
```

## Lie algebras: root system

```
julia> rs = root_system(L)
```

```
Root system of rank 2  
of type G2
```

```
julia> cartan_matrix(rs)
```

```
[ 2  -3]  
[-1   2]
```

## Lie algebras: root system

```
julia> rs = root_system(L)
Root system of rank 2
  of type G2

julia> cartan_matrix(rs)
[ 2  -3]
[-1   2]

julia> positive_roots(rs)
6-element Vector{RootSpaceElem}:
  $\alpha_1$ 
  $\alpha_2$ 
  $\alpha_1 + \alpha_2$ 
  $2\alpha_1 + \alpha_2$ 
  $3\alpha_1 + \alpha_2$ 
  $3\alpha_1 + 2\alpha_2$ 
```

# Lie algebras: Chevalley basis

```
julia> xs = chevalley_basis(L)[1]
6-element Vector{AbstractLieAlgebraElem{QQFieldElem}}:
 -3//2*b_4 - 1//2*b_8 + 3//2*b_9 + b_10 - b_11 + b_13
 -2*b_8 + b_10 - b_11 + b_13
 9//4*b_7 + 3//4*b_8 - 3//2*b_10 + 3//2*b_11 - 3//2*b_13
 27//4*b_4 + 27//4*b_5 - 9//4*b_8 - 27//4*b_9 - 9//4*b_10 + 9//4*b_11
 27//8*b_8 + 27//8*b_10 + 27//4*b_11 - 81//8*b_12 - 27//4*b_13
 81//8*b_8 - 81//4*b_10 - 81//8*b_11 + 81//8*b_13

julia> ys = chevalley_basis(L)[2]
6-element Vector{AbstractLieAlgebraElem{QQFieldElem}}:
 -2//9*b_8 + 4//9*b_10 + 2//9*b_11 - 2//3*b_12 - 2//9*b_13 + 2//3*b_14
 [...]
 -4//243*b_3 - 4//729*b_8 + 8//729*b_10 + 4//729*b_11 - 4//729*b_13

julia> hs = chevalley_basis(L)[3]
2-element Vector{AbstractLieAlgebraElem{QQFieldElem}}:
 -2*b_1 - 3*b_2
 b_1 + 2*b_2
```

Give it a try yourself!

Get in touch!

We are looking forward to your questions,  
comments, bug reports, code contributions, ...

# Practical Fast Linear Algebra Over $\mathbb{Z}$

John Abbott

---

# Practically fast linear algebra over the integers

## Joint work with Claus Fieker

### Outline:

- background & basics
- testing unimodularity
- solving linear systems (square, full-rank)
- computing (non-zero) determinants

### Existing analyses assume $A$ is *uniform*:

- most entries have similar size
- simplifies complexity analysis
- often an *unrealistic* assumption

# Basics and Background

- Work with square matrix  $A \in \text{Mat}_{n \times n}(\mathbb{Z})$ :
  - **dense**  $\rightarrow$  not many zero entries
  - **uniform**  $\rightarrow$  most entries same size as largest
  - **full-rank**  $\rightarrow \det A \neq 0$
- Write  $\|A\|_{\max}$  for largest entry in  $A$
- **SNF**  $A = U_L S U_R$  with  $S = \text{diag}(s_1, \dots, s_n)$  and  $s_j | s_{j+1}$
- **Hadamard bound** for  $|\det A|$ :
  - $H = \prod |r_i|$  where  $r_i$  is  $i$ -th row
  - $H^2 \in \mathbb{Z}$   $\leftarrow$  quick & easy to compute
  - for “random” matrices  $H$  is not too loose
- **Cramer's rule**:
  - When  $Ax = b$  then  $x_i = \det A^{(i)} / \det A$
  - so  $|\det A|$  is common denom for solution  $x$   
(best common denom is largest Smith invariant)
- **Cross-over**: boundary of regions where algm A, B, C or ... is best

# Background about CRT

**CRT** abbr. for *chinese remaindering*

**Example:** CRT method for determinant:

- For (distinct) primes  $p_1, p_2, p_3, \dots$  do  
    Compute  $D_{p_j} \leftarrow \det(\phi_{p_j}(A))$  where  $\phi_{p_j}$  is reduction  $\mathbb{Z} \rightarrow \mathbb{Z}/p_j\mathbb{Z}$
- Use CRT to recombine  $D_{p_1}, D_{p_2}, D_{p_3}, \dots$

**Stopping criterion:** Hadamard's bound, or "stability" (fast, but unverified)

**Complexity problem:** *quadratic* in size of  $\|A\|_{max}$

[each  $\phi_{p_k}(A)$  has cost linear in  $\|A\|_{max}$  and num. iterations is linear in  $\|A\|_{max}$ ]

**Solution:** reduce  $A$  modulo blocks of  $k$  primes together

→ *soft linear* complexity; needs more memory (use symmetric residues!)

Further complication: number of iterations depends on *stopping criterion*

# Testing for unimodularity

Developed from *Pauderis+Storjohann* (2012, 2013)

Unimodular matrix  $\longleftrightarrow \det(A) = \pm 1$

Two approaches:

- (1) compute  $\det(A)$  e.g. by CRT
- (2) double-plus-one lifting:  
(compute  $A^{-1}$  via cunning quadratic  $p$ -adic approximation)

Cross-over: when to use which method?

- Quick & easy to estimate accurately cost of (1)
- Cannot quickly estimate accurately cost of (2)  
(but can estimate a range)

## Testing for unimodularity (cont'd)

### Testing on random unimodular matrices

Table of ratio of double-plus-one to determinant:

n	125 bits	250 bits	500 bits	1000 bits	2000 bits
25	0.58	0.51	0.46	0.44	0.36
50	0.32	0.31	0.32	0.30	0.29
100	0.16	0.15	0.16	0.15	0.16
200	0.08	0.08	0.08	0.08	0.07

Usually quick for non-unimodular: e.g. for  $200 \times 200$  with 2000 bit entries

- $\sim 0.01s$  if not unimodular (typically)
- $\sim 4s$  if unimodular
- $\sim 50\text{--}60s$  to compute determinant (FLINT)

Q: But what is a “random unimodular matrix”?

# Solving a linear system

Task: solve  $Ax = b$  with  $A$  square, full-rank

Three approaches: (all are modular)

- (1) solve via CRT
- (2) solve via Dixon's method (linear  $p$ -adic approx; quadratic is no faster)
- (3) double-plus-one lifting: compute  $A^{-1}b$  via  $p$ -adic approx

Choice of stopping criterion: Hadamard or "stability"

Cross-over: when to use which method?

- (1) compute  $A^{-1}$  modulo each prime
- (2) compute  $A^{-1}$  modulo 1 prime; then matrix-vector prods
- (3) like (2) but cleverer/faster

# Computing the determinant

CRT algorithm: Base reference method

ABM algorithm:

- Solve  $Ax = b$  for  $x \in \mathbb{Q}^n$  for random rhs  $x \in \mathbb{Z}^n$   $\leftarrow$  “fairly fast”
- Let  $d$  be least common denominator of  $x$
- Compute  $\det(A)/d$  using CRT

Works well when  $|\det A|/d$  is small  $\longleftrightarrow$  “generic SNF”

Refined by Pauderis+Storjohann:

- Use  $d$  and numerators of  $x$  to yield triang. matrix  $H$
- Set  $A_{new} \leftarrow AH^{-1} \in \text{Mat}(\mathbb{Z})$ , and iterate
- Result is  $\det A = \det(A_{new}) \det(H)$

Better than ABM when SNF “a bit less generic”

Q: What to do when SNF “far from generic”?

## Computing the determinant (cont'd)

### New idea:

if  $d$  in ABM not too big, use P+S but with  $H$  being  $d$ -modular HNF

- Telescopes several P+S iterations into one step
- If  $d$  is max. Smith factor, we obtain  $|\det A|$  directly
- If  $d$  not too big, computing  $d$ -modular HNF is quick  
(can also “guess”  $d$  from a few CRT iterations)

Cross-over: what is “not too big”?

### Testing on “favourable” matrices

Table of ratio of new method to FLINT (standard ABM)

n	100 bits	200 bits	400 bits	800 bits
125	0.41	0.43	0.27	0.18
250	0.33	0.21	0.16	0.10
500	0.17	0.12	0.06	0.05
1000	0.14	0.14	0.06	0.03

Conclusion: det is much faster . . . . . (sometimes)

**Summing it up**

Max Horn

---

## Much more to come

- OSCAR has picked up quite some momentum!

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- ...and beyond (↪ discussion yesterday)
- We also need to “spread the word”: OSCAR ambassadors, website overhaul, merchandise, documentation sprints, OSCAR teaching material and so on (↪ the other discussion yesterday)
- We have much work ahead, not just technical





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- We meet every Wednesday, 11:30 and Friday, 15:00 on Zoom, join us: [🔗oscar-system.org/zoom](https://oscar-system.org/zoom)

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**Thank you for your attention!**